Amorphous Medium Language

Jacob Beal MIT CSAIL LSMAS Workshop, July 2005

The Big Picture

Programming large spatially distributed systems is too hard!

- Reduce complexity by programming continuous
 space and compiling for discrete agents
- Increase reuse by **functional composition** of processes













Wait a minute! Weren't we programming a farm?

We need an Abstraction Barrier!



What behavior we want from the space

How a network of agents reliably produce the behavior

- Amorphous Computing
 - [Coore 99], [Nagpal 01],
 [Kondacs 03]
- Paintable Computing [Butera 02]
- GHT [Ratnasamy et al. 02], TinyDB [Madden et al. 02]
- Regiment [Newton & Walsh 04]









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Amorphous Medium Language

- Global program \rightarrow local action \rightarrow global behavior
- Implicit distribution, coordination, communication
 - Program continuous space
 - Compile for discrete devices
 - Infrastructure supplied robust coordination primitives
- Functional composition of programs

(defun pest-management ()
 (in-components
 (dilate 2 (select-region (> (sense :bugs) 0.2)))
 (reduce-region #'max (sense :bugs) 0)))

AML: Two Compatible Models

- Global Model (continuous regions)
 - reduce-region, in-components, selectregion, dilate

- Discrete Model (agents, messages)
 - send, receive, sleep, sense, act



AML: Two 3 Compatible Models

- Global Model (continuous regions)
 - reduce-region, in-components, selectregion, dilate

- Neighborhood Model (continuous neighborhoods)
 - exposed-state@, reduce-neighbors
- Discrete Model (agents, messages)
 - send, receive, sleep, sense, act



AML: Two 3 4 Compatible Models

- Global Model (continuous regions)
 - reduce-region, in-components, selectregion, dilate
- Neighborhood AML (composable nbrhood processes, infrastructure)
- Neighborhood Model (continuous neighborhoods)
 - exposed-state@, reduce-neighbors
- Discrete Model (agents, messages)
 - send, receive, sleep, sense, act

Programming in AML



Compiling AML to N-AML

AML:

(defprocess pest-management ()
 (in-components
 (dilate 2 (select-region (> (sense :bugs) 0.2)))
 (reduce-region #'max (sense :bugs) 0)))

N-AML:

(defun pest-management ()
 (where
 (dilate 2 (lambda (place) (> (sense-bugs place) 0.2)))
 (gossiped-value #'max #'sense-bugs 0)))

• Compilation is straightforward, implementing reduce-region is difficult

Neighborhood Model



- Continuous region of space
- Each point is a separate agent
- Agents share (delayed) exposed state w. neighbors

Discrete Model

- Dozens to billions of simple, unreliable agents
- Distributed through space, communicating by local broadcast
- Agents may be added or removed



- No global services (e.g. time, naming, routing, coordinates)
- Relatively cheap power, memory, processing
- Agents are immobile*

*or slow and dense enough to run Virtual Stationary Nodes [Dolev et al. 05]

Places

- Places are Points/Agents
- State at a place:
 - Unique ID
 - Sensors (e.g. bug-detector, temperature)
 - Actuators (e.g. pesticide sprayer, LEDs)
 - Running process

Neighborhood Abstraction



- A process interacts with its neighborhood by:
 - setting the state exposed to its neighbors
 - sampling the state exposed by its neighbors

Process Execution Model



Handling Process Time



- Time is partially synchronous
 - Discrete rounds at each place
 - Different places have different clocks
- A process is an object containing state at round T
 - Round 0 specified as arguments, initial forms
 - Evolution specified as transition from T to T+1

Starting and Stopping Processes



- Processes run when installed at a place
 - Places lacking a process attempt to install from neighbor's exposed state
- Processes are tested for termination each round

Handling Interaction with Neighbors

- Exposed state = constructor args, designated slots
 - Terminated processes expose nothing
- Incremental integration between rounds
 - special neighborhood forms
 - e.g. (reduce-nbrs #'max v 0)
 - may implicitly use hidden state
 - result returned during next transition between rounds

Handling Process Composition



- A process produces a time-series of values
 - A composed process is a tree of processes containing other processes in their state
 - Compose by filling process slots (implicit or explicit)
 - Execution protocol must distribute through the tree

Local functions as processes

(sense-bugs place) accesses a sensor

(exposed-state@ place #'sense-bugs)→'(function #'sense-bugs)
(propagate@ place 'function state)→#'sense-bugs
(integrate-foreign-state@ place #'sense-bugs state)→nil [no effect]

► (execute-round@ place #'sense-bugs)→(sense-bugs place)

• (terminatep@ place #'sense-bugs)→nil

• An ordinary function is a process that ignores its neighbors and never terminates.

Describing Processes

(defnonlocal where ((f nonlocal) (target nonlocal)) (declare (termination (terminatep target))) (declare (exposing live)) (declare (integration (target (and live target)))) (with-state ((live nil)) (setf live (evaluate f)) (if live (evaluate target) nil)))

• *where* executes the *target* process only when the *f* process returns *true*

N-AML: Constructor & Class Defn

(definition of the definition (declare (termination (terminatep target))) (declare (exposing live)) (declare (integration (target (and live target)))) (with-state ((live nil)) (setf live (evaluate f)) (if live (evaluate target) nil))) (defclass where (nonlocal) ((**f** :accessor where-**f** :initarg :**f**) (target :accessor where-target :initarg :target) (live :accessor where-live :initarg :live))) (defun where (f target)) (let ((#:INST (make-instance 'where :f f :target target))) (with-slots (f target live) #:INST (setf live nil)) #:INST)))

N-AML: Execution & Termination



N-AML: Communication



Putting the pieces together...

```
(defun pest-management ()
  (where
   (dilate 2 (lambda (place) (> (sense-bugs place) 0.2)))
   (gossiped-value #'max #'sense-bugs 0)))
```

• Compound process created by functional composition of constructors

Simulated Execution



- Implemented in Allegro CommonLISP
- Runs 5000+ agent simulations

Future Work

- Optimization
- Discrete→Mica2 Motes
- Finish Global→Neighborhood
 - Update compiler
 - Improved reduce-region primitive

AML Contributions

- Abstraction barrier between **what** and **how** in large multi-agent systems
 - Computation models bridging from continuous regions of space to agents passing messages
 - Language supporting functional composition of processes
 - Primitives scalable to extremely large systems

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Thanks to Hal Abelson, Jonathan Bachrach, Gerry Sussman

Appendices

Isn't AML too expensive to use?

- Plentiful opportunities for optimization
- Communication is limited by congestion
 - Maximum density of bits, not number
- Power isn't as tight a constraint as often percieved
 - Powered networks: user-deployed (RoofNet), solar (DuraNode), embedded distribution (biological tissue)
 - Lots of research on energy storage
 - Tradeoff between control response and power usage

Using neighbor state

(defnonlocal dilate (n (source nonlocal)) (declare (termination (terminatep source))) (declare (exposing r)) (with-state ((r (1+ n))) (merge-nbrs (r) (setf r (min r (nbr r)))) (when (evaluate source) (setf r -1)) (incf r) (<= r n)))</pre>

• The *dilate* process returns *true* within *n* units of points where the *source* process returns *true*

Using neighbor state

(defnonlocal gossiped-value (fuser (source nonlocal) &optional (base unspecified)) (declare (exposing value)) (declare (termination nil)) (with-state ((value base)) (setf value (funcall fuser (evaluate source) (reduce-nbrs fuser value base)))))

• The *gossiped-value* process uses *fuser* to reduce the values of the *source* process to the same summary value at every place